

## HOW LARGE CAN THE COEFFICIENTS OF A POWER SERIES BE?

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*Dedicated to Professor János Galambos on the occasion of his 70th birthday*

**Abstract.** The paper is inspired by the problem of estimating the deviation of two discrete probability distributions in terms of the supremum distance between their generating functions over the interval  $[0, 1]$ . Under certain conditions on the tail it is clarified how large can the terms of a real sequence be if the sup norm of its generating function is known.

### 1. Introduction

Let  $A_1, \dots, A_n$  be an arbitrary collection of events in an arbitrary probability space. Let  $N$  denote the number of events that occur. In many cases we have to determine the distribution of the random variable  $N$ , or, at least, to estimate the probability that none of the events occur. Such problems typically arise when stochastic methods are applied in combinatorics, see [1].

The probability  $P(N = 0)$  can be estimated in several ways. In messy situations, where the dependence structure of the events is rather complicated, sieve methods,

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like the Rényi sieve, can sometimes help. Those methods provide Bonferroni type lower or upper bounds of the form

$$(1.1) \quad P(N = 0) \leq (\geq) \sum_M c(M) P\left(\bigcap_{i \in M} A_i\right),$$

where in the sum  $M$  runs over the subsets of  $\{1, 2, \dots, n\}$ , and the  $c(M)$  are real constants. The interested reader is referred to the excellent monograph by Galambos and Simonelli [3].

Such inequalities can easily be transformed into bounds for the probability generating function of  $N$ . By [4, Theorem 1], together with (1.1) we also have

$$(1.2) \quad g_N(x) = E(x^N) \leq (\geq) \sum_M c(M) P\left(\bigcap_{i \in M} A_i\right) (1-x)^{|M|},$$

for  $0 \leq x \leq 1$ . If we want to show the asymptotic Poissonity of  $N$ , in the way above we can estimate the difference between  $g_N(x)$  and the generating function of the corresponding Poisson distribution (i.e., that with expectation equal to  $EN$ ). Now the question is: how to estimate the difference of probabilities, if we have bounds for the difference of generating functions?

This problem can be reformulated in the following way. Let  $\mathcal{F}$  be the set of real power series  $f(x) = \sum_{k=0}^{\infty} a_k x^k$  such that  $\sum_{k=0}^{\infty} |a_k| \leq 2$  and  $\sum_{k=0}^{\infty} a_k = 0$ . When we have two discrete probability distributions  $\mathbf{p} = (p_0, p_1, \dots)$  and  $\mathbf{q} = (q_0, q_1, \dots)$  with generating functions  $g_{\mathbf{p}}(x)$  and  $g_{\mathbf{q}}(x)$ , resp., then  $f = g_{\mathbf{p}} - g_{\mathbf{q}} \in \mathcal{F}$ , and  $a_k = p_k - q_k$ . We want to estimate  $a_k$  in terms of  $\Delta = \max_{0 \leq x \leq 1} |f(x)|$ .

The main difficulty of the problem is in the restriction that we only know  $f$  over the real interval  $[0, 1]$ , not in a whole neighbourhood of the origin on the complex plane.

The first results in this direction appeared in [4]. It is shown there that the coefficients can not be estimated uniformly, and that for every  $\ell = 0, 1, \dots$  and  $\varepsilon > 0$  there exists a constant  $C$ , depending on  $\ell$  and  $\varepsilon$ , such that

$$|a_{\ell}| \leq C \Delta^{1-\varepsilon}.$$

This was improved in [5] to

$$|a_{\ell}| \leq \Delta \exp\left(2\left(\ell \log \frac{1}{\Delta}\right)^{4/5}\right).$$

On the other hand, the limitations of such estimations are illustrated by the following counterexamples, borrowed from [5].

Let  $\mathbf{p} = (p_0, p_1, \dots)$  be a fixed discrete probability distribution such that  $p_k > 0$  for every  $k = 0, 1, \dots$ , and

$$(1.3) \quad \limsup_{k \rightarrow \infty} \frac{1}{2k} \log \frac{1}{p_k} < \infty.$$

Let  $\ell$  be a positive integer and  $C$  a sufficiently small positive constant. Then for every sufficiently small positive  $\Delta$  there exists a discrete probability distribution  $\mathbf{q} = (q_0, q_1, \dots)$ , such that  $\max_{0 \leq x \leq 1} |g_{\mathbf{p}}(x) - g_{\mathbf{q}}(x)| = \Delta$ , and

$$(1.4) \quad |p_\ell - q_\ell| > C \Delta \left( \log \frac{1}{\Delta} \right)^{2\ell}.$$

Analogous result holds for probability distributions  $\mathbf{p}$  having a tail that is lighter than exponential.

Suppose that, instead of (1.3), we have

$$(1.5) \quad \limsup_{k \rightarrow \infty} \frac{1}{h(k)} \log \frac{1}{p_k} = v,$$

where  $v$  is positive and finite,  $h$  is a positive, continuous, increasing function, regularly varying at infinity with exponent  $\alpha$ , and  $\lim_{k \rightarrow \infty} h(k)/k = \infty$  (hence  $\alpha \geq 1$ ). Let  $\ell$  be a positive integer and  $C$  a sufficiently small positive constant. Then for every sufficiently small positive  $\Delta$  there exists a discrete probability distribution  $\mathbf{q}$ , such that  $\max_{0 \leq x \leq 1} |g_{\mathbf{p}}(x) - g_{\mathbf{q}}(x)| = \Delta$ , and

$$(1.6) \quad |p_\ell - q_\ell| > C \Delta \left( h^{-1} \left( \log \frac{1}{\Delta} \right) \right)^{2\ell}.$$

These examples inspired our results in Section 2. We are going to drop the condition  $\sum_{k=0}^{\infty} |a_k| \leq 2$ , but in that case (for  $\ell > 0$ )  $|a_\ell|$  can be arbitrary large, no matter how small  $\Delta$  is, see Theorem 2.

In order to derive upper bounds in the form of the right hand sides of (1.4) and (1.6) we have to impose additional conditions on the sequence  $(a_k)$  in consideration.

## 2. Results

Let us start with a fundamental lemma.

The following theorem is a variant of a result by V. A. Markov, who proved a similar theorem on the extremal properties of Chebyshev polynomials over the interval  $[-1, 1]$  (see Chapter 2 of [2]).

**Theorem 1.** *Consider an arbitrary polynomial of the form  $Q_n(x) = \sum_{k=0}^n a_k x^k$ . Introduce  $\Delta = \max_{0 \leq x \leq 1} |Q_n(x)|$ . Then*

$$(2.1) \quad |a_k| \leq \frac{n}{k+n} \binom{k+n}{2k} 2^{2k} \Delta.$$

For  $k > 0$  equality holds if and only if  $Q_n(x) = \pm \Delta T_n(2x - 1)$ , where  $T_n$  is the degree  $n$  Chebyshev polynomial of the first kind, defined as  $T_n(\cos \theta) = \cos(n\theta)$ .

Let us remark that

$$(2.2) \quad \frac{n}{k+n} \binom{k+n}{2k} 2^{2k} \leq \frac{(2n)^{2k}}{(2k)!}$$

by the inequality of arithmetic and geometric means, and the ratio of the two sides tends to 1 as  $k$  is fixed and  $n \rightarrow \infty$ .

**Proof.** We may assume that  $a_k = 1$ . Then  $T_n(2x_i - 1) = (-1)^i$  for  $x_i = \cos^2 \frac{i\pi}{2n}$ ,  $i = 0, 1, \dots, n$ . Let  $T_n(2x - 1) = \sum_{k=0}^n d_k x^k$ . Suppose

$$|d_k| \Delta < 1 = \max_{0 \leq x \leq 1} |T_n(2x - 1)|.$$

Let  $p(x) = T_n(2x - 1) - d_k Q_n(x)$ , then  $p(x_i)$  is positive or negative, according as  $i$  is even or odd. Thus  $p(x)$  has  $n$  distinct roots in the interval  $(0, 1)$ . Let them be denoted by  $y_1, y_2, \dots, y_n$ , then

$$p(x) = (d_n - d_k a_n) \prod_{i=1}^n (x - y_i),$$

hence the coefficient of  $x^k$  in  $p(x)$  is equal to

$$(-1)^{n-k} (d_n - d_k a_n) \sum_{1 \leq i_1 < \dots < i_{n-k} \leq n} y_{i_1} \dots y_{i_{n-k}}.$$

It should be 0, which is a contradiction. Consequently,  $|a_k| = 1 \leq |d_k| \Delta$ .

Suppose  $k > 0$  and  $a_k = \Delta d_k$ . Let  $p(x) = \Delta T_n(2x - 1) - Q_n(x)$ , and suppose that  $p$  is not identically equal to 0. Then  $p(x_i) \geq 0$  for even values of  $i$ , and  $p(x_i) \leq 0$  for odd  $i$ . Hence  $p$  has a zero in every closed interval  $[x_i, x_{i-1}]$ ,  $i = 1, 2, \dots, n$ . If  $p(x_i) = 0$  for some  $1 < i < n$ , it must be a multiple root, since  $p$  does not change sign at  $x_i$ . This shows that there are exactly  $n$  zeros in  $[0, 1]$ , if each root is counted up to its multiplicity. In addition, if  $x_n = 0$  is a root, it must be single. Again, let  $y_1, y_2, \dots, y_n$  be the roots, and consider the coefficient of  $x^k$  in  $p(x)$ ,

$$0 = (-1)^{n-k} (\Delta d_n - a_n) \sum_{1 \leq i_1 < \dots < i_{n-k} \leq n} y_{i_1} \dots y_{i_{n-k}}.$$

Since  $k > 0$ , the sum on the right-hand side must have at least one positive term, leading to a contradiction.

Finally, it is known [5] that

$$d_k = (-1)^{n-k} \frac{n}{k+n} \binom{k+n}{2k} 2^{2k}.$$

This can be shown by induction on  $n$ , basing on the recursion formula  $T_n(2x - 1) = 2(2x - 1)T_{n-1}(2x - 1) - T_{n-2}(2x - 1)$ . ■

Let us add that for  $k = 0$  equality can also hold for polynomials other than  $Q_n(x) = \pm \Delta T_n(2x - 1)$ . For example, every polynomial of the form  $Q_n(x) = \pm \Delta(1 - xP_{n-1}(x))$  will do, where  $\deg P_{n-1} \leq n - 1$ , and  $0 \leq P_{n-1}(x) \leq 2$  on  $[0, 1]$ .

In what follows we consider real sequences  $(a_k)$  such that the series  $\sum_{k=0}^{\infty} a_k$  is absolutely convergent. As in Section 1, we introduce the generating function

$$(2.3) \quad f(x) = \sum_{k=0}^{\infty} a_k x^k, \quad 0 \leq x \leq 1,$$

and the notation

$$(2.4) \quad \Delta = \sup_{0 \leq x \leq 1} |f(x)|.$$

First we point out that  $a_\ell$  cannot be estimated without any further restriction.

**Theorem 2.** *For arbitrary  $\ell > 0$  and  $\Delta > 0$  we have*

$$\sup \left\{ |a_\ell| : \max_{0 \leq x \leq 1} |f(x)| = \Delta \right\} = \infty.$$

**Proof.** Let  $f(x) = \Delta T_n(2x - 1)$ , then  $|a_\ell| = \Delta \cdot 2^{2\ell} \frac{n}{\ell + n} \binom{\ell + n}{2\ell}$ , which tends to infinity with  $n$ . ■

This is the reason why we set additional conditions. We formulate them in the flavor of (1.4) and (1.5).

**Theorem 3.** *Let  $h$  be a positive function defined on  $[0, \infty)$ , such that  $\frac{h(x)}{x}$  tends nondecreasingly to a limit  $\varrho$  as  $x \rightarrow \infty$ ;  $0 < \varrho \leq \infty$ . Suppose  $\Delta \leq e^{-h(\ell)}$ , and*

$$(2.5) \quad \sum_{k=n+1}^{\infty} |a_k| \leq K e^{-h(n)}, \quad n \in \mathbb{N},$$

with some positive constant  $K$ . Then

$$(2.6) \quad |a_\ell| \leq (K + 1) C_\ell \Delta \left[ h^{-1} \left( \log \frac{1}{\Delta} \right) \right]^{2\ell},$$

where  $C_\ell = \frac{2^{2\ell}}{(2\ell)!}$ . On the other hand, for every positive  $K' < \left(\frac{\varrho}{2+\varrho}\right)^{2\ell}$  and every sufficiently small  $\Delta > 0$  there exists a sequence  $(a_k)$  such that (2.4) and (2.5) hold, and

$$(2.7) \quad |a_\ell| \geq K' C_\ell \Delta \left[ h^{-1} \left( \log \frac{1}{\Delta} \right) \right]^{2\ell}.$$

**Theorem 4.** *Suppose the conditions of Theorem 3 are met except that*

$$(2.5') \quad \sum_{k=0}^{\infty} |a_k| e^{h(k)} \leq K < \infty$$

is satisfied instead of (2.5). Then (2.6) follows. On the other hand, for every positive  $K' < \left(\frac{\varrho}{2+\varrho}\right)^{2\ell}$  and every sufficiently small  $\Delta > 0$  there exists a sequence  $(a_k)$  such that (2.4), (2.5'), and (2.7) hold.

**Proof of Theorems 3 and 4.** Note that condition (2.5') implies (2.5), for

$$\sum_{k=n+1}^{\infty} |a_k| \leq e^{-h(n)} \sum_{k=0}^{\infty} |a_k| e^{h(k)}.$$

Suppose (2.5) holds. Choose

$$n = \left\lceil h^{-1} \left( \log \frac{1}{\Delta} \right) \right\rceil,$$

then  $e^{-h(n)} \leq \Delta$ . Cutting the power series (2.3) into two at the  $n$ th term we obtain that

$$\sup_{[0,1]} \left| \sum_{k=0}^n a_k x^k \right| \leq \Delta + \sum_{k>n} |a_k| \leq \Delta + K e^{-h(n)} \leq (K+1)\Delta.$$

Hence Theorem 1 and inequality (2.2) immediately imply (2.6).

For the other direction choose  $n$  so that

$$2n + h(n) \leq \log \frac{K}{\Delta} < 2(n+1) + h(n+1),$$

and let  $f = Q_n \Delta$ , where  $Q_n(x) = T_n(2x-1)$ . Then (2.4) is fulfilled. For the coefficient of  $x^\ell$  in  $f(x)$  we have

$$|a_\ell| = \Delta \cdot 2^{2\ell} \frac{n}{\ell+n} \binom{\ell+n}{2\ell} \sim C_\ell \Delta n^{2\ell}.$$

Firstly, suppose that  $\varrho < \infty$ . Then  $h^{-1}(x) \sim x/\varrho$  as  $x \rightarrow \infty$ , and

$$(2 + \varrho)n \sim 2n + h(n) \sim \log \frac{1}{\Delta}$$

as  $\Delta \rightarrow 0$ , that is, as  $n \rightarrow \infty$ . Hence

$$n \sim \frac{1}{2 + \varrho} \log \frac{1}{\Delta} \sim \frac{\varrho}{2 + \varrho} h^{-1}\left(\log \frac{1}{\Delta}\right) \sim \frac{\varrho}{2 + \varrho} \left[ h^{-1}\left(\log \frac{1}{\Delta}\right) \right].$$

Secondly, let  $\varrho = \infty$ . Then  $h^{-1}(x)/x$  converges nonincreasingly to 0. It follows that  $h^{-1}(h(n) + c) \sim n$  if  $c = c(n) = o(h(n))$ . Indeed, for  $c \geq 0$  we have

$$n = h^{-1}(h(n)) \leq h^{-1}(h(n) + c) \leq \frac{h^{-1}(h(n))}{h(n)} (h(n) + c) = n \left(1 + \frac{c}{h(n)}\right) \sim n.$$

Similarly, for  $c < 0$  all inequalities hold reversed. Since

$$h(n) + 2n - \log K \leq \log \frac{1}{\Delta} \leq h(n + 1) + 2(n + 1) - \log K,$$

we obtain that

$$h^{-1}(h(n) + 2n - \log K) \leq h^{-1}\left(\log \frac{1}{\Delta}\right) \leq h^{-1}(h(n + 1) + 2(n + 1) - \log K),$$

thus

$$n \sim h^{-1}\left(\log \frac{1}{\Delta}\right) \sim \left[ h^{-1}\left(\log \frac{1}{\Delta}\right) \right].$$

All we have left is to show that condition (2.5') is satisfied.

$$\sum_{k=m+1}^n |a_k| e^{h(k)} \leq \sum_{k=m+1}^n \frac{(2n)^{2k}}{(2k)!} \Delta e^{h(k)} \leq \Delta e^{h(n)+2n} \leq K. \quad \blacksquare$$

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